# New Operating System Version 1.0

# For the DSE VZ200/300 Computer's

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#### Introduction

#### What are the aims of this Project?

The main aim of this project is to allow easier access to any peripheral devices attached to the VZ by any piece of software already written for the VZ with little or no modification.

The main assumptions before you commence this project are:-

- 1) You have reasonable knowledge of Z80 Assembly Language Programming.
- 2) You own a VZ with at least a 16K RAM Expansion.
- 3) You have time to do the work required.

#### The Vision

A VZ Computer with a maximum of 4 Megabytes of memory (available to all programs - even BASIC Programs), Multiple Disk Drives attached (including a 720K 3.5 inch Drive), Enhanced Keyboard, optional RAM Disk Software, A Real-Time Clock keeping track of the Date and Time allowing Alarm Features etc, a Mouse or other pointing device, an 80 column display for better viewing of data, and to round off the set-up - a Sound Blaster Card allowing you to record and play back sounds from various input devices (TV, Video, Radio, Stereo etc).

We are aiming High!

#### The BASIC's

All components in **NOS** are designed to be replaceable, so that you may configure your system in any way you would like. The system has been designed to emulate the DOS System of the popular IBM Personal Computers. That is, it will provide similar functionality but will **not** neccessarily be compatible. **NOS** will consist of nine basic components:-

B:	BOOT.EXE	This is the NOS Initialisation Program.
B:	MEMORY	This is the Memory Manager and will look after all memory allocation regardless of its location.
B:	PROCMAN	This is the Process Manager and looks after the loading/executing of all Programs.
B:	SCRNMAN	This is the Screen/Display Manager and looks after the different Graphics Modes etc.
B:	DISKMAN	This is the Disk Manager and shields all applications from the native Disk Drive Commands.
B:	PRINTMAN	This is the Print Manager and looks after the translation of control codes and provides a Printer Spooler.
T:	CONFIG	This file contains the Peripheral Device Driver installation/configuration information.
B:	COMMAND	This is the Command processor and contains the basic reserved words allowed at the command line.
T:	AUTOEXEC	This file contains any other Programs the User may wish to load when starting <b>NOS</b> .

#### **Notes/Assumptions:**

- 1) The files listed above must all live on the same diskette which will from this point forth be referred to as the NOS Boot Disk.
- 2) The **NOS Boot Disk** is a standard VZ Diskette and as such can only be read from a VZ Disk Drive. This means that we will have two limitations: 80K Maximum Boot Disk size and also we cannot boot from a different type of Disk Drive yet. These are not major hurdles and will be resolved in the future.

## The Initialisation Process (BOOT.EXE)

As mentioned above, the entry point into **NOS** will be via **BOOT.EXE** which, as with standard VZ disk software, will be started with the usual BRUN"BOOT.EXE". Once loaded, the Boot program will perform some minor tasks such as setting up internal pointers and reserving permanent memory areas for itself. It will then commence to load the other programs from the boot disk in the order listed above. The loading order is essential as some of the processes reply upon the previous components being there to operate correctly. Below is the pseudo code for the Boot process, see **Appendix A** for the executable version.

 $Begin Program\ `Boot'$ 

Origin: 40960; Origin = A000H

Disable Interrupts ; Prevent the System from having problems

; whilst initialisation occurs.

Reserve 25 Bytes for the **System Information Frame (SIF)** 

at locations 31465 to 31489

; See later for further details

Initialise SIF: ; Values are for NOS 1.0 only.

Store NOS Version Description Pointer;=31490Store High Order NOS Version;=1Store Low Order NOS Version;=0Initialise the System Flags;=0Initialise the Max Process ID;=255

Set the Base System Stack Pointer ; = 31597 + 40 (Default Stack Size)

Set the Top System Stack Pointer ; = 31597

Set the Current System Stack Pointer ; = Base System Stack Pointer

Store IM1 Table Pointer ;= 31507Initialise IM1 Function Count ;= 0Store the Critical Error Handler Address ;= 31636

; All following addresses are known to the Boot Program only. Any Process wishing to use the following

; information should go via the SIF to retain compatibility with future versions of NOS.

Store Version Description String at location 31490 (16 Bytes) ; "NOS Version 1.0"

Initialise Interrupt Device Table at location 31507 ; 31507 - 31596 = 0

(max 30 Functions - 3 Bytes each = 90 Bytes)

Reserve 40 Bytes for the NOS System Stack at location 31597

Install the Critical Error Handler at location 31636 ; See later for further details

Install the Interrupt Processing Subroutine for IM1 ; See later for further details

executed via an RST 38H instruction.

; Load the remaining NOS System Components

EXECUTE[ MEMORY ] ; Install the Memory Manager
EXECUTE[ PROCMAN ] ; Install the Process Manager
EXECUTE[ SCRNMAN ] ; Install the Screen Manager
EXECUTE[ DISKMAN ] ; Install the Disk Manager
EXECUTE[ PRINTMAN ] ; Install the Printer Manager

Call SYS Interpreter [ CONFIG ] ; Read the CONFIG file line by line and

; perform the specified commands

EXECUTE[ COMMAND ] ; Install the Command Line Processor

Call Batch Interpreter[ AUTOEXEC ] ; Read the Autoexec file line by line and

; perform the specified commands.

Enable Interrupts

End Program 'Boot' ; Return control to the User (NOS Installed)

## The System Information Frame (SIF)

This is an area of memory which contains all the important parameters (offsets) needed by the system to handle the multi-processing and the communication between processes and **NOS**. The definition of **SIF** will continue to change as we add further information as **NOS** expands. Once an offset has been declared in **SIF**, it will not be removed. Also, the location of **SIF** must never relocate and is fixed at location 31465 decimal (7AE9H). This way backwards compatibility will be retained so that software written for **NOS** version 1.0 will still run under a future version without change.

The current definition of SIF is as follows:-

Offset	# of Bytes	Description		
0	2	Size of SIF		
2	2	NOS Version Description Pointer		
4	1	High Order NOS Version		
5	1	Low Order NOS Version		
6	2	Base System Stack Pointer		
8	2	Top System Stack Pointer		
10	2	Current System Stack Pointer		
12	2	Reserved		
14	1	Reserved		
15	2	IM1 Interrupt Device Table		
17	1	No. entries in IM1 Interrupt Device Table		
18	2	Reserved		
20	1	Reserved		
21	1	System Flags:		
		Bit 0 = Stack In Use Indicator 0 - SP is the Current Process Stack 1 - SP is the System Stack  Bits 1 & 2 = Interrupt Processing Indicator 00 - No Interrupt is being Processed 01 - Reserved 10 - IM1 Interrupt is being Processed		
		Bits 3 & 4 = Stack Request after Interrupt 00 - No Request 01 - Request System Stack after Interrupt 10 - Request Current Process Stack after Interrupt 11 - Reserved  Bits 5 - 7 = Reserved for future expansion Set to Zero (000)		
22	1	Last Process ID. (Total No. of Processes currently loaded)		
23	2	Address of the Critical Error Handler		

### **The Critical Error Handler**

The Critical Error Handler is just that, a centralised error checking routine which given an error code produces a message and then either returns control to allow the calling routine to cleanup as best it can, or it terminates the parent process if a major error has occured. For a full list of error messages refer to **Appendix**??.

## The IM1 Interrupt Processing Subroutine

The standard VZ hardware configuration is such that only the IM1 type interrupts of the Z80 CPU have been enabled. This is the simplest of the three types of interrupts the Z80 permits. When an interrupt occurs an **RST 38H** instruction is executed which in turn calls a small service routine. Currently, the only interrupting device on a standard VZ is the Graphics Controller Chip which interrupts 50 times every second. The current interrupt service routine only looks after screen updates, such as flashing the cursor, displaying messages and looking after the background colours.

Over the years a lot of software and hardware enhancements have meant that the standard service routine was replaced with a customised device specific version. This was a good solution until the time came where more than one add-on was required at any one time. The Interrupt Service routine installed by **NOS** will attempt to overcome this problem.

For the initial version of **NOS** we will only be considering IM1 type interrupts. It is most likely that this will change in a future version, but hardware enhancements will also be required.

In **NOS** there is a table which contains a list of functions, in priority order, which are to be called every time an interrupt occurs. The address of this table is stored at offset 15 of **SIF** and is known as the IM1 Device Table. For this version of NOS there is a limit of 30 functions which can be called every time an interrupt occurs. Each function looks after a specific device and whenever called polls the device to see if any events occurred since the last interrupt. The one major drawback to this method is that if we do not poll a device often enough we will loose information. The only way to resolve this is via some form of hardware enhancement which is beyond the scope of the initial version.

Questions:
How do we handle interrupts of devices which require different polling priorities

there is a limit on the amount of time we can spend servicing interrupts before the system slows down.

### The Memory Manager

The Memory Manager controls the allocation/distribution of memory for use by all programs. Memory on the VZ can be defined as one of four types:-

- 1. System Memory
- 2. Extended Memory
- 3. Expanded Memory
- 4. Screen Memory

All but the last type of memory is looked after by the Memory Manager. The minimum expected configuration is System memory only, although this will severely restrict the user and negate any of the benefits of NOS.

#### System Memory

This is the area of memory between 7AE9H and BFFFH, and is mandatory. Applications should not use this memory unless absolutely necessary. This memory is used by NOS for loading Device Drivers, lookup tables, file control blocks etc.

To access System Memory, the following primitives are available:

i) Allocate\_System

- This marks a block of system memory as being used and returns a pointer to it.

ii) ReAllocate\_System iii) Release\_System

- This increases the size of an allocated block of system memory and returns a pointer to it. - This releases a block of allocated system memory.

iv) TotalFree\_System

- The returns the amount of unused system memory.

v) LargestFree\_System vi) SmallestFree\_System

- The returns the size of the largest unused block of system memory. - The returns the size of the smallest unused block of system memory.

vii) FirstFree\_System viii) NextFree\_System

- The returns the size of the first unused block of system memory. - The returns the size of the next unused block of system memory.

#### Extended Memory

This is the area of memory starting at C000H. Depending upon the hardware configuration, the end address can vary. The memory in this region may also be bank switched.

To access extended memory, the following primitives have been provided:

Allocate\_Global i) ii) ReAllocate\_Global iii) Release\_Global

iv) TotalFree\_Global

- This marks a block of extended memory as being used and returns a pointer to it.

- This increases the size of an allocated block of extended memory and returns a pointer to it. - This releases a block of allocated extended memory. - This returns the amount of unused extended memory.

v) LargestFree\_Global vi) SmallestFree Global vii) FirstFree\_Global

- This returns the size of the largest unused block of extended memory. - This returns the size of the smallest unused block of extended memory. - This returns the size of the first unused block of extended memory.

viii) NextFree\_Global - This returns the size of the next unused block of extended memory. ix) NoBanks\_Global - This returns the number of extended memory banks.

x) TotalMem\_Global - This returns the total amount of extended memory.

Please note that these primitives return a three byte (24 bit) pointer. The following additional primitives have been provided to store and retrieve data from extended memory:

StoreByte\_Global

- This stores a byte of data in global extended memory.

ii) ReadByte\_Global

- This reads a byte of data from global extended memory.

iii) StoreBlock\_Global

- This transfers a block of data to global extended memory.

iv) ReadBlock\_Global

- This reads a block of data from global extended memory.

v) TransferBlock\_Global

- This transfers a block of data from one area of global extended memory to another.

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The Process Manager	

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# The Process Information Frame (PIF)

The current definition of PIF is as follows:-

Offset	# of Bytes	Description
0	2	Size of PIF
2	2	Start Address
4	1	Bank No
5	1	Process ID
6	2	Base Process Stack Pointer
8	2	Top Process Stack Pointer
10	2	Current Process Stack Pointer
12	2	Function Table Pointer
14	1	Number of External Functions
15	1	unication Call/Results  Called Process No
16	1	Sub Function No
17	1	ES: Call Error Status/Single Return Value
18	2	RV1: Double Byte Return Value
20	2	RV2: Double Byte Return Value

# **Appendix ? - Error Messages produced by the Critical Error Handler**

Error Code	Terminate/Suppress	Reason for Error	Error Message
0	-	No Error	-
1	Suppress	Invalid Process ID	A Requested Process is not available
2	Suppress	Invalid Function Call	A Requested Function is not available
3	Terminate	Stack Underflow Error	A Stack Error has Occurred.Application is terminating
4	Terminate	Stack Overflow Error	A Stack Error has Occurred.Application is terminating

NB: The second column indicates that the parent process can suppress these error messages and resolve the problem programatically.

# Appendix ? - Standard Process Identifier's

Process ID	Description
; NOS Services	
0	Reserved
1	Memory Manager
2	Process Manager
3	Screen Manager (Video I/O)
4	Disk Manager
5	Print Manager
; Device Drivers	s - Standard Hardware
6	Keyboard Input/Output
7	Video Input/Output
8	Cassette Input/Output
9	Speaker Output
10 11	Parallel Printer Output Joystick Input
12	Disk Input/Output
; Device Drivers	s - Non-standard Hardware Extensions Support
13	Mouse Input
14	Real-Time Clock Input/Output
15	Serial Input/Output
; Device Drivers	s - Logical Devices
16	Standard Input
17	Standard Output
18	Standard Printer
19	Standard Error
;Reserved	
20 - 24	Reserved for Future Expansion.
	ger allocates from 25 to 255

# Appendix ? - Useful Definitions

To make the code easier to read, the following definitions have be created:

SIF	EOU	31465
SYSTEM_FLAGS	EQU	SIF + 19
PROCESS_LOOKUP_TABLE	EQU	SIF + 23
SYSTEM_STACK	EQU	SIF + 10
LAST PROCESS ID	EOU	SIF + 22

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## Appendix ?? - The IM1 Interrupt Handler

**Upon Entry:** Stack contains all registers except IY & IX.

A, (SYSTEM\_FLAGS) **INTR:** LD ; Indicate that we are servicing IM1 Interrupts. AND 11111101b ; Preserve all bits other than bits 1 & 2. (SYSTEM\_FLAGS), A LD POP HL; Remove the return address to the ROM routine from the stack. PUSH ΙY ; Save the Current Process Frame Pointer. **PUSH** ; Save the IX Register. BIT 0,A; Check if System Stack is already in use. JR NZ, INT1 LD IY, SIF ; Set Frame Pointer to SIF. INT1: HL, -20 LD ; Recalculate the Current Stack Pointer allowing for the ADD HL, SP ; registers already placed there by the ROM routine. LD (IY + 10), L; HL = Current Stack Pointer, could be the Process Stack. LD (IY + 11), H; SP is unaffected. VSTACK ; Validate that the stack integrity is OK. CALL SERVICE INTERRUPTS ; Refer Appendix ?? for further Information. POP ΙX ; Restore the IX Register. POP ΙY ; Restore the Current Process Frame Pointer. POP HL; Restore registers placed on the stack by the ROM routine. POP DE POP LD A, (SYSTEM\_FLAGS) ; Reload the System flags. BIT ; Test for System Stack request. 3, A NZ, INT3 JR BIT 4, A ; Test for Process Stack request. JR NZ, INT4 INT2: AND 11100001b ; Reset: Interrupt Indicator & Stack request. LD (SYSTEM\_FLAGS), A POP AF EI RETI INT3: BIT ; Test if System Stack already in use NZ, INT2 JR ; Goto INT2 if so. AND 11100001b ; Reset: Interrupt Indicator & Stack request. SET 0, A ; Indicate SP = System Stack. LD (SYSTEM\_FLAGS), A POP AF ; Restore Registers placed there by ROM routine.

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	; Because we are messing around with the stack pointer we cannot rely upon it.			
	LD POP	(SAVE_REG1), HL HL	; Save the current value of HL; HL = Commencement Address when returning from interruging; This is not the return address to the ROM routine as this has already been removed.	
	LD PUSH LD EI RETI	SP, (SIF + 10) HL HL, (SAVE_REG1)	; SP = Current System Stack Pointer. ; Restore the Commencement Address ; Restore value of HL	
INT4:	BIT JR	0, A Z, INT2	; Test if System Stack in use ; Goto INT2 if not.	
	AND RES LD	11100001b 0, A (SYSTEM_FLAGS), A	; Reset: Interrupt Indicator & Stack request. ; Indicate SP = Current Process Stack.	
	POP	AF	; Restore Registers placed there by ROM routine.	
	; Becau	se we are messing around w	ith the stack pointer we cannot rely upon it.	
	LD	(SAVE_REG1), HL	; Save the current value of HL	
	POP	HL	; HL = Commencement Address when returning from interru ; This is not the return address to the ROM routine as this has ; already been removed.	
	LD	(SAVE_REG2), HL	; Save it	
	LD LD	L, (IY + 10) H, (IY + 11)	; HL = Current Process Stack Pointer	
	LD LD	SP, HL HL, (SAVE_REG2)	; SP = Current Process Stack Pointer	
	PUSH LD EI RETI	HL HL, (SAVE_REG1)	; Restore Commencement Address ; Restore value of HL	

#### $; These \ storage \ are as \ are \ known \ only \ to \ the \ Interrupt \ Service \ Routine.$

SAVE\_REG1: DEFS 2 SAVE\_REG2: DEFS 2

## **Appendix ?? - Stack Validation Routine**

**Upon Entry:** HL = Stack Pointer.

**Upon Exit:** HL register altered,

> Flags register altered, All other registers preverved.

PUSH DE **VSTACK:** 

> E, (IY + 6)LD

LD D, (IY + 7)EX

; DE = Base of Stack. DE, HL ; DE = Current Stack Pointer. ; HL = Base of Stack.

OR ; Clear carry flag.

HL. DE SBC

JR ?, VSTACK\_UFLOW ; Stack Underflow Error.

LD L, (IY + 8)

LD ; HL = Top of Stack.  $H,\,(IY+9)$ 

OR

EX DE, HL ; HL = Current Stack Pointer.

; DE = Top of Stack.

; Restore the A register.

SBC HL, DE

?, VSTACK\_OFLOW ; Stack Overflow Error. JR

OR ; Clear carry flag.

POP DE

RET

VSTACK\_UFLOW: LD E, A ; Preserve the A register.

LD A, 3

CALL CRITICAL\_ERRORHANDLER

LD A, E

POP DE

RET

VSTACK\_OFLOW: LD E, A ; Preserve the A register.

LD A, 4

CALL CRITICAL\_ERRORHANDLER

LD A, E ; Restore the A register.

POP DE

RET

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## **Appendix ?? - Inter-Process Function Call Routine (RST 30H)**

**Upon Entry:** IY = Current Process **PIF**.

**Upon Exit:** IY = Current Process **PIF**.

All other registers destroyed.

IP\_CALL: LD A, (SYSTEM\_FLAGS) ; Get the current System Flags.

> BIT ; Is the call occurring during interrupt processing.

; Goto IP\_CALL1 if so. NZ, IP\_CALL1 JR

; Wait for an interrupt. This serves two purposes: HALT

; i) Yields to O/S allowing background

processes to occur.

; ii) Ensures stack integrity.

IP\_CALL1: LD H, 0

L, (IY + 13)LD ; HL = Called Process No.

A, (LAST\_PROCESS\_ID) LD ; A = ID of last Process loaded. CP ; Check if Called Process No is valid

LD ; Set Errorflag. A, 1

JΡ ?, CRITICAL\_ERRORHANDLER ; "Invalid Process ID".

ADD HL, HL ; \*2

PUSH HL

; \*2 (\*4) ADD HL, HL ADD HL, HL ; \*2 (\*8)

POP DE

ADD HL, DE (\*10)ADD HL, HL ; \*2 (\*20)

EX ; DE = Table Offset DE, HL

LD IX, (TABLE\_BASE\_ADDRESS)

ADD ; IX = Pointer to Called Process **PIF**. IX, DE

LD E, (IX + 10)

LD D, (IX + 11); DE = Function Table Pointer.

LD A, D ; Check that the Process ID is valid.

OR

Z, CRITICAL\_ERRORHANDLER

LD A, (IX + 12)A = No of available functions.

L, (IY + 14)LD ; L = Function No.

CP ; Compare them. L

LD A, 2

?,  ${\bf CRITICAL\_ERRORHANDLER}$  ; "Invalid Function Call". JP

XOR Α ; A = 0, carry flag cleared.

SLA L

RLA

LD H, A

ADD HL, DE ; HL = Table Address of function.

LD A, (SYSTEM\_FLAGS)

; Check if System stack in use. BIT 0, A NZ, IP\_CALL3 ; Goto IP\_CALL3 if not. JR

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	; System	Stack is not in use.	
	BIT JR	1, A Z, IP_CALL2	; Check if interrupt being processed. ; Goto IP_CALL2 if not.
	; Current	tly Processing Interrupts.	
	EX	DE, HL	; Save Table Address in DE
	LD ADD	HL, 0 HL, SP	; HL = Current Stack Pointer
	LD LD	(IY + 8), L (IY + 9), H	; Save current Stack pointer.
	CALL	VSTACK	; Validate that the stack integrity is OK.
	EX LD	DE, HL SP, (SYSTEM_STACK)	; HL = Table address of Function. ; Set the Stack Pointer to the System Stack.
	SET LD JR	0, A (SYSTEM_FLAGS), A IP_CALL3	; Indicate System Stack in use. ; GOTO IP_CALL3
IP_CALL2:	SET LD HALT	2, A (SYSTEM_FLAGS), A	; Request System stack on next interrupt. ; Wait for an interrupt.
IP_CALL3:	LD CALL	A, (IX + 2) BANK_SWITCH	; A = Bank of Called Process. ; = OUT (7FH), A
	LD INC LD	E, (HL) HL D, (HL)	; DE = Function Address.
	PUSH EX	IX (SP), IY	; Save Called Process <b>PIF</b> . ; Exchange it with current Process <b>PIF</b> .
	LD BIT JR	A, (SYSTEM_FLAGS) 1, A Z, IP_CALL4	; Check if interrupt being processed. ; Goto IP_CALL4 if not.
	; Current	tly Processing Interrupts.	
	SET LD	0, A (SYSTEM_FLAGS), A	; Indicate Process Stack in use.
	LD LD	L, (IY + 8) H, (IY + 9)	; HL = Called Process Stack Pointer.
	LD	SP, HL	; SP = Called Process Stack Pointer.
	JR	IP_CALL5	
IP_CALL4:	SET LD	3, A (SYSTEM_FLAGS), A	; Request current Process Stack on next interrupt.
	HALT		; Wait for an interrupt.
IP_CALL5:	LD PUSH EX JP	HL, (CLEANUP_ADDRESS) HL DE, HL (HL)	; Execute Function.

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