New Operating System Version 1.0

For the DSE VZ200/300 Computer's

Contents

Introduction

What are the aims of this Project?

The main aim of this project is to allow easier access to any peripheral devices attached to the VZ by any piece of software already written for the VZ with little or no modification.

The main assumptions before you commence this project are:-

- 1) You have **reasonable** knowledge of Z80 Assembly Language Programming.
- 2) You own a VZ with at least a 16K RAM Expansion.
- 3) You have time to do the work required.

The Vision

A VZ Computer with a maximum of 4 Megabytes of memory (available to all programs - even BASIC Programs), Multiple Disk Drives attached (including a 720K 3.5 inch Drive), Enhanced Keyboard, optional RAM Disk Software, A Real-Time Clock keeping track of the Date and Time allowing Alarm Features etc, a Mouse or other pointing device, an 80 column display for better viewing of data, and to round off the set-up - a Sound Blaster Card allowing you to record and play back sounds from various input devices (TV, Video, Radio, Stereo etc).

We are aiming High!

The BASIC's

All components in **NOS** are designed to be replaceable, so that you may configure your system in any way you would like. The system has been designed to emulate the DOS System of the popular IBM Personal Computers. That is, it will provide similar functionality but will **not** neccessarily be compatible. **NOS** will consist of nine basic components:-

Notes/Assumptions:

- 1) The files listed above **must** all live on the same diskette which will from this point forth be referred to as the **NOS Boot Disk.**
- 2) The **NOS Boot Disk** is a standard VZ Diskette and as such can only be read from a VZ Disk Drive. This means that we will have two limitations: 80K Maximum Boot Disk size and also we cannot boot from a different type of Disk Drive yet. These are not major hurdles and will be resolved in the future.

The Initialisation Process (BOOT.EXE)

As mentioned above, the entry point into **NOS** will be via **BOOT.EXE** which, as with standard VZ disk software, will be started with the usual BRUN"BOOT.EXE". Once loaded, the Boot program will perform some minor tasks such as setting up internal pointers and reserving permanent memory areas for itself. It will then commence to load the other programs from the boot disk in the order listed above. The loading order is essential as some of the processes reply upon the previous components being there to operate correctly. Below is the pseudo code for the Boot process, see **Appendix A** for the executable version.

 ; All following addresses are known to the Boot Program only. Any Process wishing to use the following ; information should go via the **SIF** to retain compatibility with future versions of **NOS**.

The System Information Frame (SIF)

This is an area of memory which contains all the important parameters (offsets) needed by the system to handle the multi-processing and the communication between processes and **NOS**. The definition of **SIF** will continue to change as we add further information as **NOS** expands. Once an offset has been declared in **SIF**, it will not be removed. Also, the location of **SIF** must never relocate and is fixed at location 31465 decimal (7AE9H). This way backwards compatibility will be retained so that software written for **NOS** version 1.0 will still run under a future version without change.

The current definition of SIF is as follows:-

The Critical Error Handler

The Critical Error Handler is just that, a centralised error checking routine which given an error code produces a message and then either returns control to allow the calling routine to cleanup as best it can, or it terminates the parent process if a major error has occured. For a full list of error messages refer to **Appendix ??**.

The IM1 Interrupt Processing Subroutine

The standard VZ hardware configuration is such that only the IM1 type interrupts of the Z80 CPU have been enabled. This is the simplest of the three types of interrupts the Z80 permits. When an interrupt occurs an **RST 38H** instruction is executed which in turn calls a small service routine. Currently, the only interrupting device on a standard VZ is the Graphics Controller Chip which interrupts 50 times every second. The current interrupt service routine only looks after screen updates, such as flashing the cursor, displaying messages and looking after the background colours.

Over the years a lot of software and hardware enhancements have meant that the standard service routine was replaced with a customised device specific version. This was a good solution until the time came where more than one add-on was required at any one time. The Interrupt Service routine installed by **NOS** will attempt to overcome this problem.

For the initial version of **NOS** we will only be considering IM1 type interrupts. It is most likely that this will change in a future version, but hardware enhancements will also be required.

In **NOS** there is a table which contains a list of functions, in priority order, which are to be called every time an interrupt ocurrs. The address of this table is stored at offset 15 of **SIF** and is known as the IM1 Device Table. For this version of NOS there is a limit of 30 functions which can be called every time an interrupt occurs. Each function looks after a specific device and whenever called polls the device to see if any events occurred since the last interrupt. The one major drawback to this method is that if we do not poll a device often enough we will loose information. The only way to resolve this is via some form of hardware enhancement which is beyond the scope of the initial version.

Questions:

How do we handle interrupts of devices which require different polling priorities there is a limit on the amount of time we can spend servicing interrupts before the system slows down.

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The Memory Manager

The Memory Manager controls the allocation/distribution of memory for use by all programs. Memory on the VZ can be defined as one of four types:-

- 1. System Memory
- 2. Extended Memory
- 3. Expanded Memory
- 4. Screen Memory

All but the last type of memory is looked after by the Memory Manager. The minimum expected configuration is System memory only, although this will severely restrict the user and negate any of the benefits of **NOS**.

System Memory

This is the area of memory between 7AE9H and BFFFH, and is mandatory. Applications should not use this memory unless absolutely necessary. This memory is used by NOS for loading Device Drivers, lookup tables, file control blocks etc.

To access System Memory, the following primitives are available:

Extended Memory

This is the area of memory starting at C000H. Depending upon the hardware configuration, the end address can vary. The memory in this region may also be bank switched.

To access extended memory, the following primitives have been provided:

- i) Allocate_Global This marks a block of extended memory as being used and returns a pointer to it.
- ii) ReAllocate_Global This increases the size of an allocated block of extended memory and returns a pointer to it.
- iii) Release_Global This releases a block of allocated extended memory.
iv) TotalFree_Global This returns the amount of unused extended memory.
	- This returns the amount of unused extended memory.
- v) LargestFree_Global This returns the size of the largest unused block of extended memory.
- vi) SmallestFree_Global This returns the size of the smallest unused block of extended memory.
- vii) FirstFree_Global This returns the size of the first unused block of extended memory.
	-
- viii) NextFree_Global This returns the size of the next unused block of extended memory.
ix) NoBanks_Global This returns the number of extended memory banks. - This returns the number of extended memory banks.
- x) TotalMem_Global This returns the total amount of extended memory.

Please note that these primitives return a three byte (24 bit) pointer. The following additional primitives have been provided to store and retrieve data from extended memory:

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- i) StoreByte_Global This stores a byte of data in global extended memory.
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- ii) ReadByte_Global This reads a byte of data from global extended memory.
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-
-
- iii) StoreBlock_Global This transfers a block of data to global extended memory.
iv) ReadBlock_Global This reads a block of data from global extended memory.
- iv) ReadBlock_Global This reads a block of data from global extended memory.
v) TransferBlock_Global This transfers a block of data from one area of global exte - This transfers a block of data from one area of global extended memory to another.

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The Process Manager

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The Process Information Frame (PIF)

The current definition of PIF is as follows:-

Appendix ? - Error Messages produced by the Critical Error Handler

NB: The second column indicates that the parent process can suppress these error messages and resolve the problem programatically.

Appendix ? - Standard Process Identifier's

Appendix ? - Useful Definitions

To make the code easier to read, the following definitions have be created:

Appendix ?? - The IM1 Interrupt Handler

Upon Entry: Stack contains all registers except IY & IX.

Page 16 New Operating System 1.0 Appendix ?? - Inter-Process Function Call Routine (RST 30H) Upon Entry: IY = Current Process **PIF**. **Upon Exit:** IY = Current Process **PIF**. All other registers **destroyed**. **IP_CALL:** LD A, (SYSTEM_FLAGS) ; Get the current System Flags. BIT 1, A ; Is the call occurring during interrupt processing. JR NZ, IP_CALL1 ; Goto IP_CALL1 if so. HALT ; Wait for an interrupt. This serves two purposes: ; ; i) Yields to O/S allowing background processes to occur. ; ii) Ensures stack integrity. **IP_CALL1:** LD H, 0
LD L, $(IV + 13)$; HL = Called Process No. LD A, (LAST_PROCESS_ID) ; A = ID of last Process loaded. CP L
 $\begin{array}{ccc}\n\text{CPP} & \text{L} \\
\text{D} & \text{A. 1}\n\end{array}$; Check if Called Process No is valid

: Set Errorflag. A, 1 ; Set Errorflag. JP ?, **CRITICAL_ERRORHANDLER** ; "Invalid Process ID". ADD HL, HL *2 PUSH HL
ADD HL, ADD HL, HL $; *2$ (*4)
ADD HL, HL $; *2$ (*8) HL, HL POP DE ADD HL, DE ; $(^{*}10)$
ADD HL, HL ; *2 $(^{*}20)$ ADD HL, HL ; *2 (*20) EX DE, HL ; DE = Table Offset LD IX, (TABLE_BASE_ADDRESS)
ADD IX, DE IX , DE ; IX = Pointer to Called Process **PIF**. LD E, $(IX + 10)$
LD D, $(IX + 11)$ \mathcal{L} DE = Function Table Pointer. LD A, D ; Check that the Process ID is valid. OR E
JP Z. **Z, CRITICAL_ERRORHANDLER** LD A, $(IX + 12)$; A = No of available functions.
LD L, $(IV + 14)$; L = Function No. ; $L =$ Function No. CP L ; Compare them. LD $A, 2$

JP $2, C1$?, **CRITICAL_ERRORHANDLER** ; "Invalid Function Call". XOR A ; $A = 0$, carry flag cleared. SLA L RLA $H. A$ ADD HL, DE ; HL = Table Address of function. LD A, (SYSTEM_FLAGS)
BIT 0, A BIT 0, A ; Check if System stack in use.

JR NZ, IP_CALL3 ; Goto IP_CALL3 if not.

; Goto IP_CALL3 if not.

